Highly Scalable Metadata Search and Indexing

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Challenge: scalable search mechanisms

- Indexing is a critical issue
 - Speed and effectiveness of search limit the usability of very large scale storage systems
- Very large scale indexes are often resource-intensive
 - Google and Yahoo have web-scale indexes, but they use thousands of processors to do it!
 - Performance is high (memory resident indexes)
- Indexing can take advantage of locality
 - Users typically aren't searching over the whole file system
 - Users may not have permissions to see everything
- Challenges
 - Building indexes that scale
 - Building less resource-intensive indexes
 - Building indexes that leverage locality
 - Incorporating security into indexes
 - Integrity: failure of a centralized index can be difficult to recover from







Challenge: gathering metadata for indexes

- Indexes are only as good as the information that goes into them
- Critical types of metadata include
 - Content
 - Domain-specific techniques for gathering it
 - May need domain-specific search mechanisms
 - Provenance
 - How was the data generated?
 - On what data and programs does this file depend?
- Challenges
 - How can provenance be tracked efficiently?
 - How can domain-specific metadata be handled?
 - Gathered?
 - Indexed as part of the file system?







Challenge: data mining in mass storage

- Large storage systems contain a lot of useful data
 - Can be difficult to fully utilize
- Traditional data mining techniques may not be effective
 - Infeasible to read out the entire storage system for data mining
- Two potential approaches
 - Index the data when it's written to storage, and use the indexes for mining
 - Distribute computation to the storage devices, allowing them to run in parallel





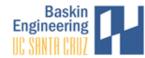


Our approach:partitioned indexes

- Break up single large index into many smaller indexes
 - Each subindex covers a manageable amount of data
 - Individual subindexes can be searched quickly
 - Individual subindexes can be rebuilt after corruption (or to reflect incremental changes)
- ◆ Problem: can't search all subindexes for every query
 - No better than what we do today (and maybe worse)
- Problem: subindex search needs to be efficient
 - Can't cache all subindexes in memory
- Leverage locality?

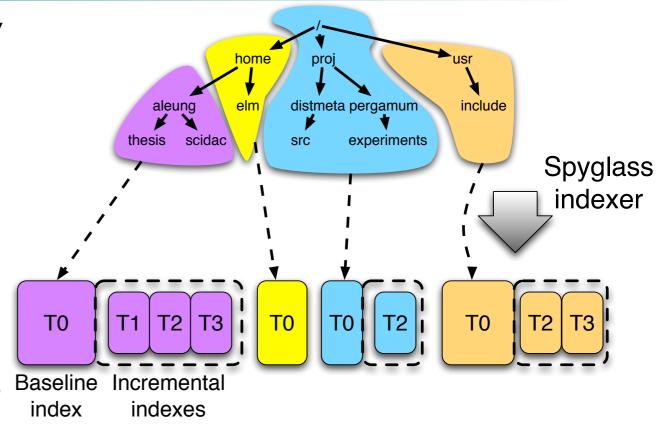


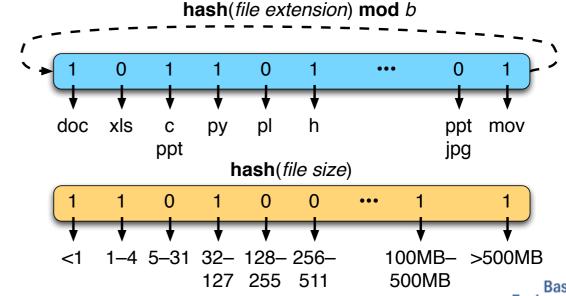




Spyglass design

- Partition file system hierarchy by subtree
 - Each subtree is an independent subindex
- Summarize contents of each subindex
 - Quickly rule out entire subindexes that can't satisfy the query
- Log incremental changes
 - Rebuild index when there are "enough" changes
- Integrity is much easier
 - Rebuild subindex, not entire index









Current status

- Prototype implemented for attributes
 - File size, type, owner, etc.
 - Content being done now...
- Used k-d trees for individual indexes
 - Straightforward to optimize tree structure for each individual subtree
 - Recently-accessed subtrees cached in memory
 - Only takes a few milliseconds to read and search a tree
- Implementation tested using 300 million files crawled at a "large storage company"
 - Locality really helps!
- Performance is very good
 - Compared against standard databases
 - Queries on Spyglass were 10–6000 times faster!







Ongoing research

- Extending attribute-based search to content
 - Will locality help as much?
 - Will compression help to reduce the size of the indexes?
- Scaling indexes to trillions of files
 - Existing systems (databases or otherwise) can't handle such large systems without massive amounts of hardware
 - Will summarization work with millions of indexes?
- Non-hierarchical file systems
 - Use similarity (or something else) to group files?
 - Talk on this earlier today
- Extending indexing to archives





